

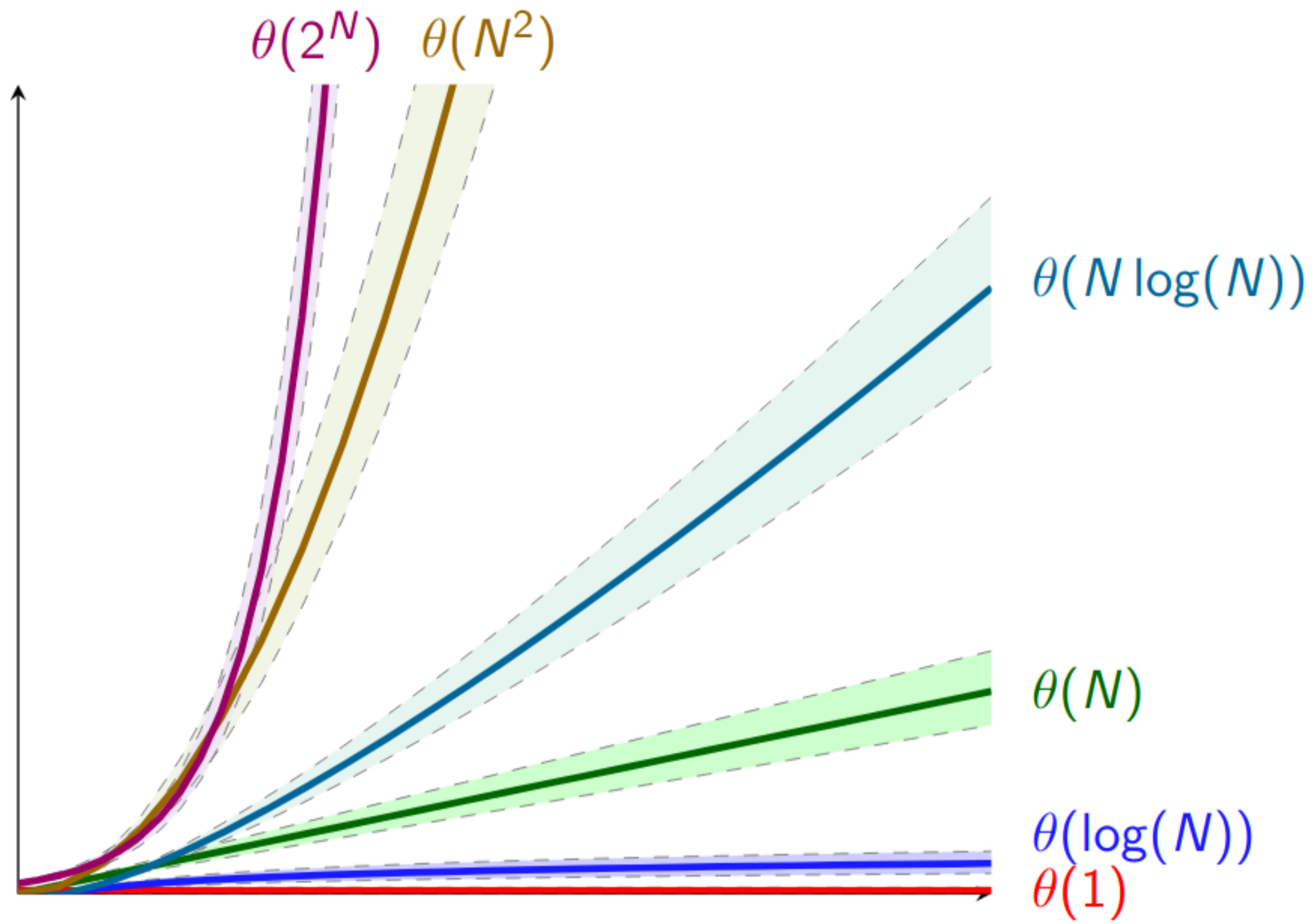
CSE 410 – Advanced Data Structures

Topic 03: Background

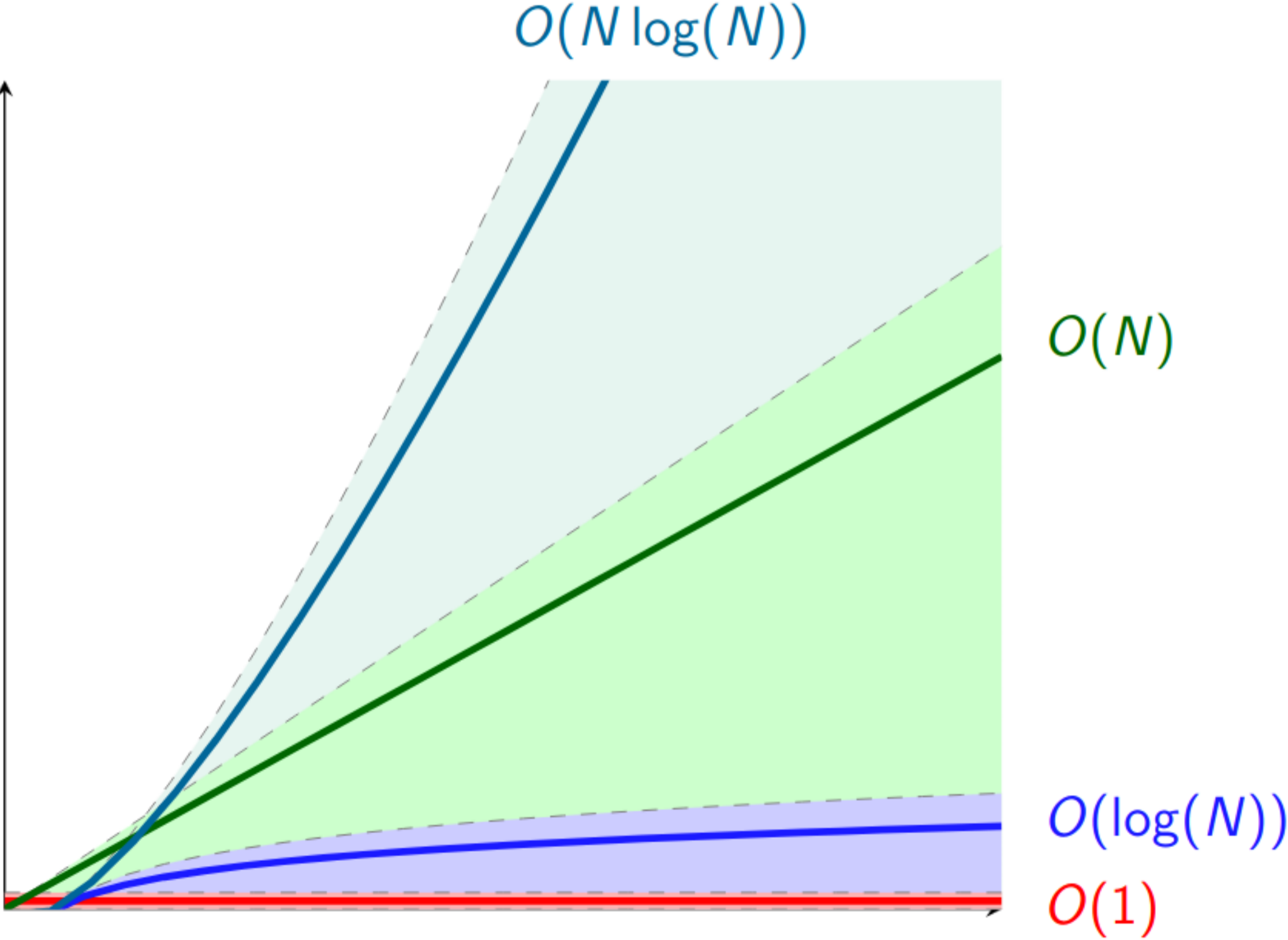
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Complexity Classes



Big-O Bounds



Analyzing Code

```
File: DFS.java
1 public Map<Vertex,Vertex> traverse(Graph graph, Vertex start, TodoList todos)
2 {
3     Map<Vertex,Vertex> backEdges = new HashMap<>() ← Q1)
4     todos.add( new Edge(null, start) ) ← A
5     while(!todos.isEmpty()) ← E
6     {
7         Edge nextEdge = todos.remove() ← R
8         if(!backEdges.contains(nextEdge.to)) ← O(1) Expected, O(N) general
9         {
10            backEdges.put(nextEdge.to, nextEdge.from) ← O(1) Expected, O(N) general
11            for(edge : graph.getOutEdges(nextVertex)) ← O(deg(v))
12            {
13                todos.add(edge) ← A
14            }
15        }
16    }
17 }
```

each vertex once

$O(\text{deg}(v))$ times

^{add} ^{remove}
A, E, R are parameters of TodoList
#E, #V are parameters of the graph

Amortized Bounds

File: AmortizedQueue.java

```
1 class AmortizedQueue<T>
2 {
3     SinglyLinkedList<T> head = new SinglyLinkedList();
4     SinglyLinkedList<T> tail = new SinglyLinkedList();
5
6     public void enqueue(T value)
7     {
8         tail.prepend(value) —  $O(1)$ 
9     }
10    public T dequeue()
11    {
12        if(head.isEmpty){
13            head = tail.reverse() —  $O(N)$ 
14            tail = new SinglyLinkedList()  $O(1)$ 
15        }
16        return head.removeFirst() —  $O(1)$ 
17    }
18 }
```

Dequeue is $O(N)$
in general, but...

usually $O(1)$

Each $O(N)$ run
"credits" us for
 N calls to dequeue

The total runtime for N calls to dequeue is Always $O(N)$
Amortized $O(1)$

Expectations

Unqualified Runtime:

Always guaranteed

Amortized Runtime

Guaranteed average over N calls

→ individual calls may be slower

Expected Runtime

Statistically Likely

Collection ADTs

Collection, Elements Identified by Position

List (Sequence)

Collection, Elements Identified by Position in Sort Order

Sorted List

Collection, Elements Identified by Their Value

Set

Collection, Elements Identified by Key

Value

Data Structure Building Blocks

List

↳ Array, Linked list, Vector (aka Array List, Array, Buffer)

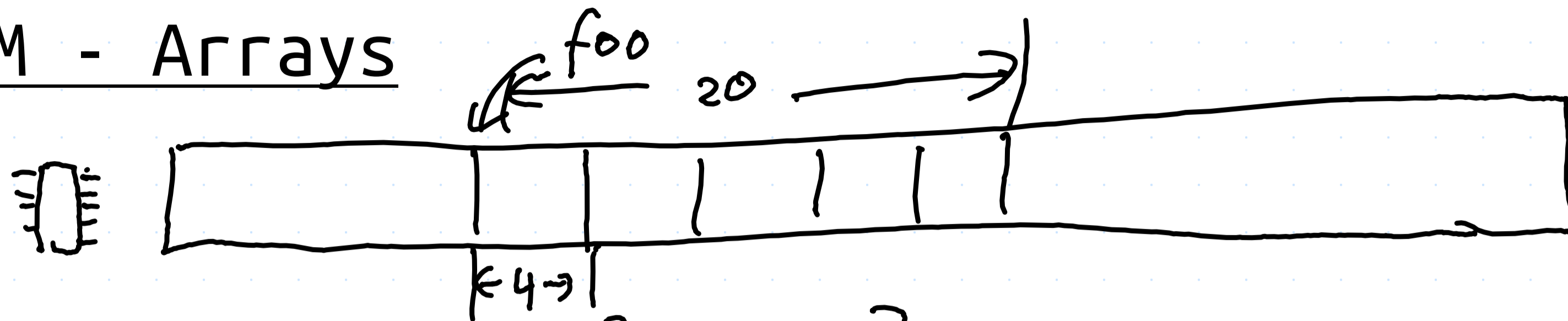
Set

↳ Tree, Hash Table

Map

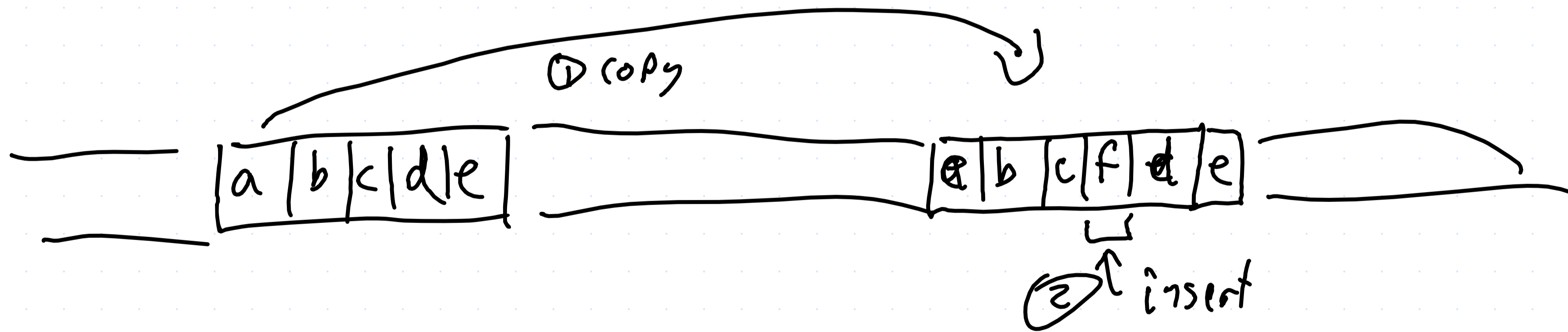
↳ Tree, Hash Table

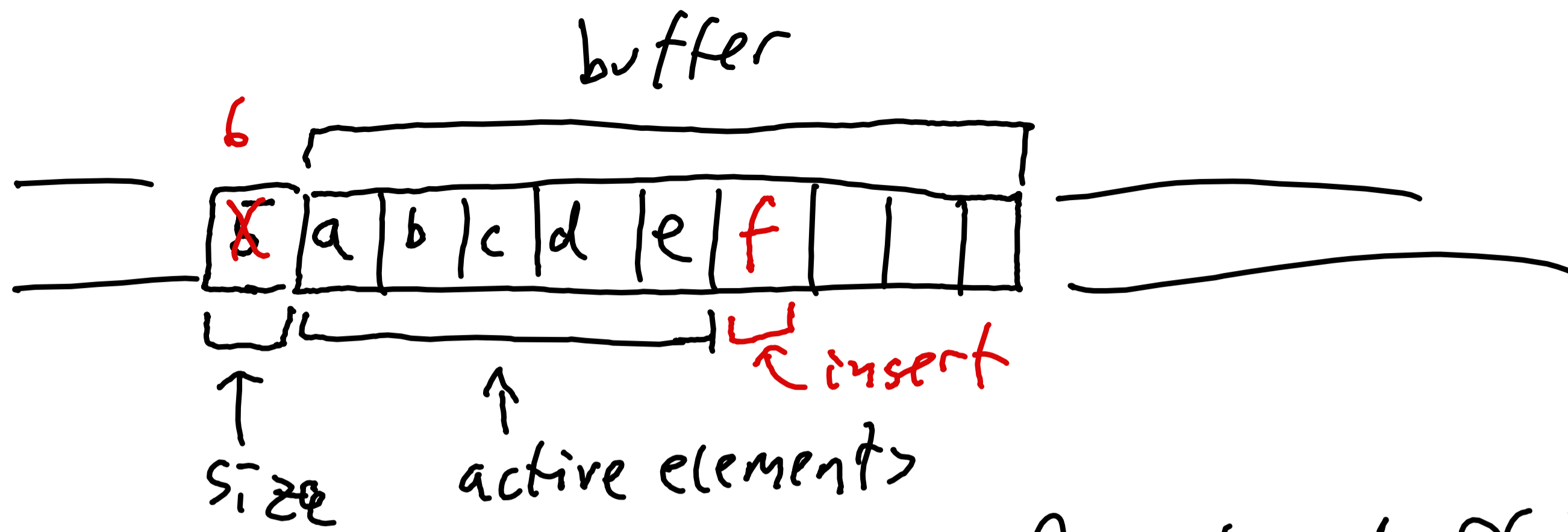
RAM - Arrays



let foo: [u32; 5] = _____
 ↑ 4b = 32/8
 ↑

$foo[3] \leftarrow foo + 3 \cdot \underbrace{4}_{\text{size of one element}}$

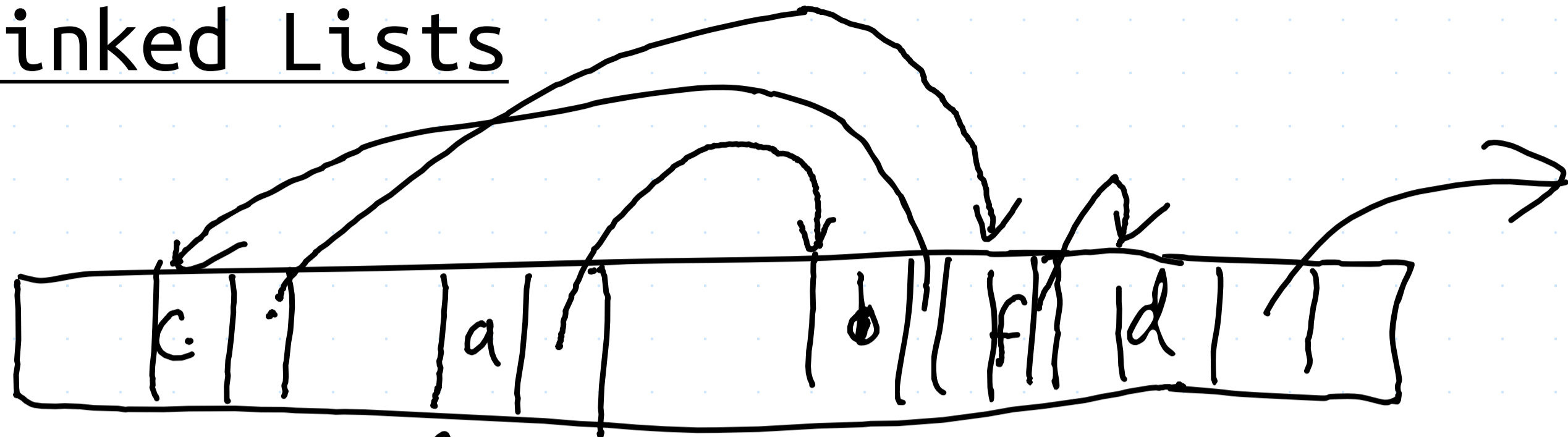




Amortized $O(1)$ insert

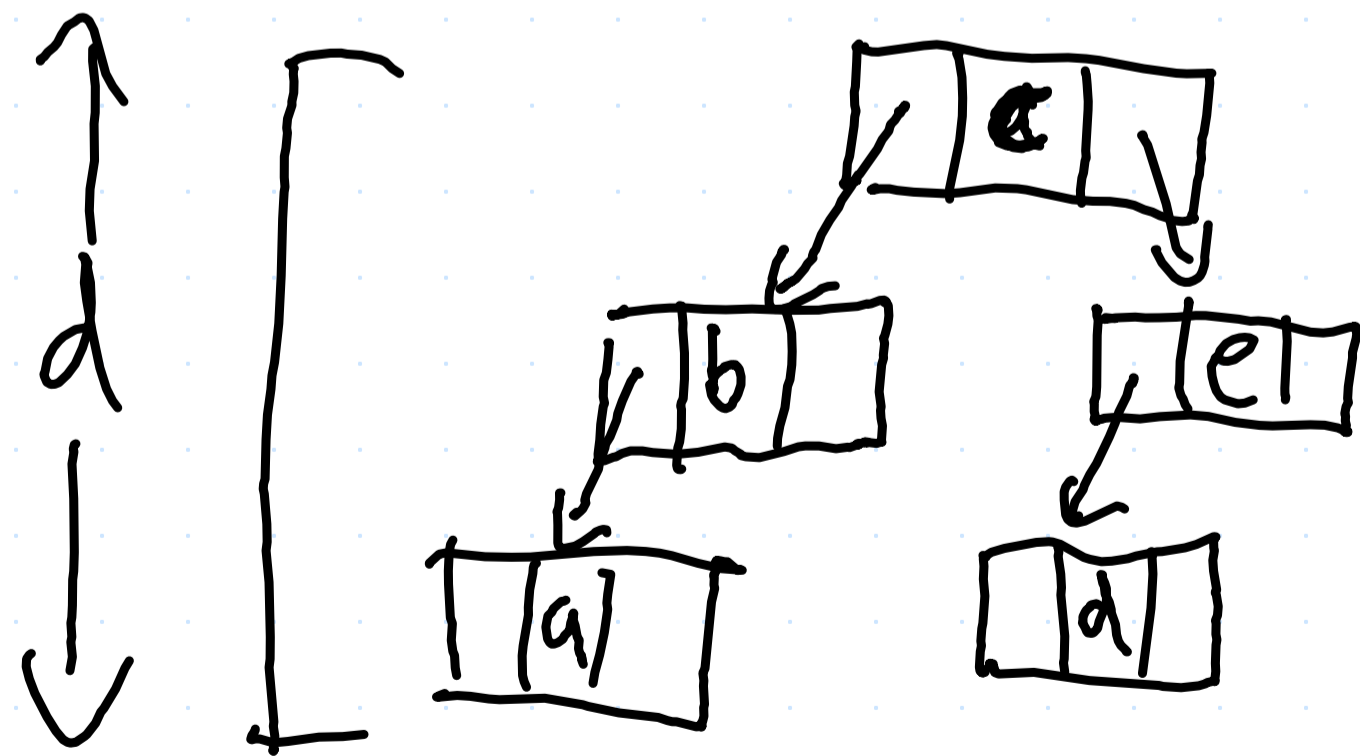
- ArrayBuffer (Scala)
- ArrayList (Java)
- Vector (Vec) (Rust)

RAM - Linked Lists



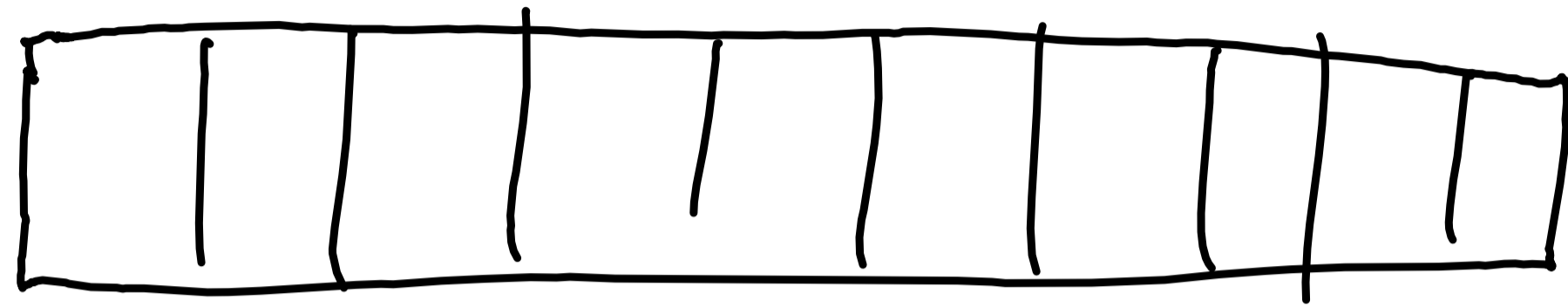
a, b, c, d, e

$O(i)$ to find i th element



Binary search tree
 $O(d) \approx O(N)$

Balanced Binary Search tree
 $\hookrightarrow d = O(\log N)$



↑ $h(x) = 5$

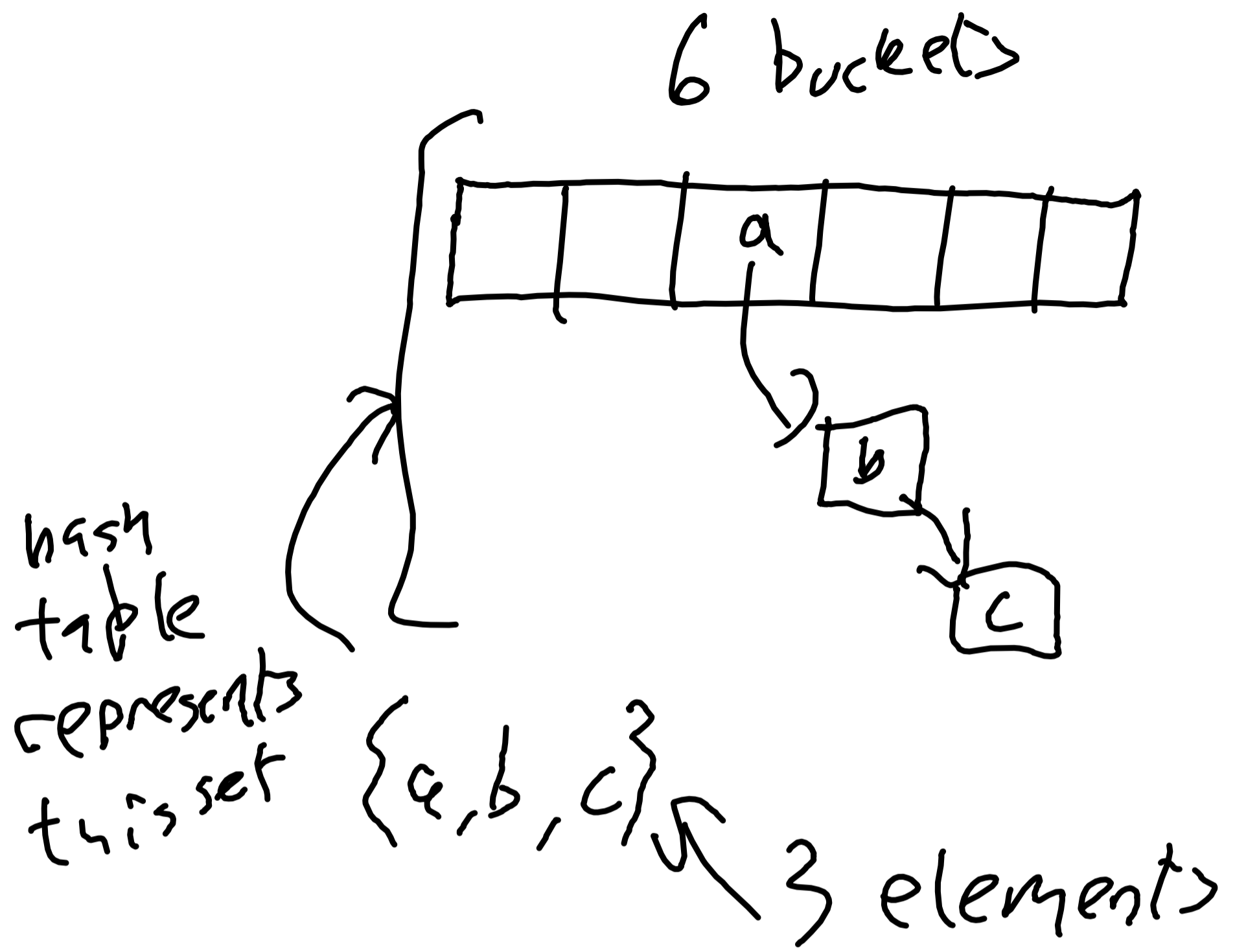
Hash function

↳ Produces a pseudo random output

↳ Output is deterministic

$$h(42) = h(42)$$

$h(42) \% N \rightarrow$ position in an N element array
where 42 is stored



$$h(a) = h(b) = h(c)$$

Hash Table w/ Chaining

↳ Expected $O(1)$

More Structure in the Data

Primitive: Int, Float, Char, ^{usually} {String}

Combine primitives as

↳ Struct / Tuple } - fixed set of fields
- Types are independent
- fields identified by name

↳ Collection } - variable number of elements
↳ list/vector - All of the same type
↳ Set/Map - Identified by position (list)
- by key (map)
- Properties
↳ Uniqueness → Set

Collection of
Struct of
Primitive 1, Prim 2, ...

Table, Dataframe, Relation

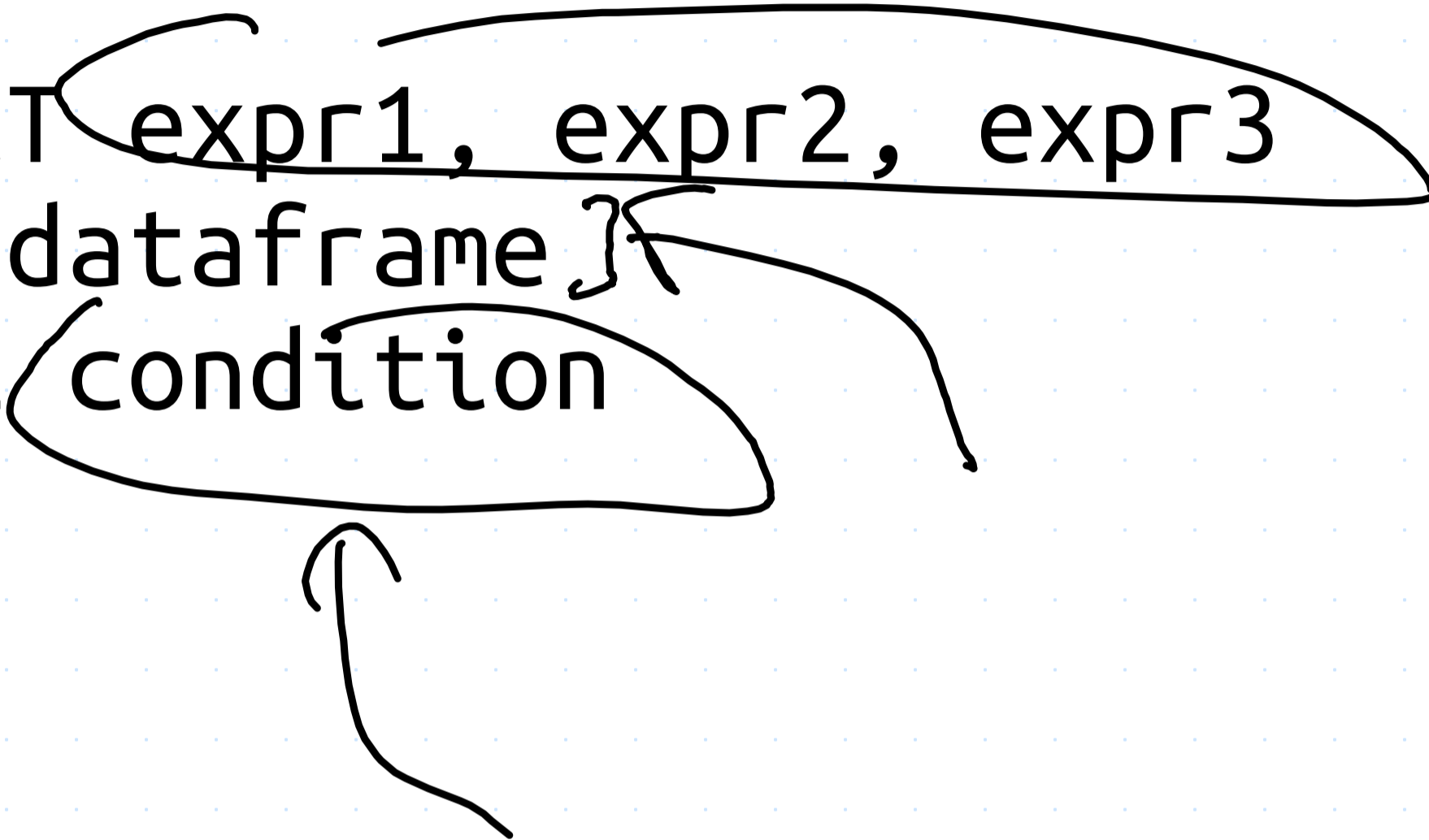
CartId	Account
1	alice
2	bob
3	carol

CartId	Item
1	apple
1	banana

SQL

```
SELECT expr1, expr2, expr3  
FROM dataframe  
WHERE condition
```

← for each record



$$d(v) + A + \sum_{v \in V} \left[E + R + d(v) + \sum_{e \text{ adjacent to } v} A \right]$$

$$= d(v) + A + \sum_{v \in V} \left[E + R + d(v) + A \cdot \text{deg}(v) \right]$$

$$= d(v) + A + \sum_{v \in V} \left[E + R + A \cdot \text{deg}(v) \right]$$

$$= d(v) + A + \sum_{v \in V} [E + R] + \sum_{v \in V} [A \cdot \text{deg}(v)]$$

$$= d(v) + A + \#V [E + R] + A \cdot \#E$$

$$= \#V \cdot (E + R) + A \cdot \#E$$


```
SELECT DISTINCT expr1, expr2, expr3  
FROM dataframe  
WHERE condition
```

Traversal Algo	Todo List	E	R	A	Traversal Runtime
BFS	Queue	$O(1)$	$O(1)$	$O(1)$	$\#V + \#E$
DFS	Stack	$O(1)$	$O(1)$	$O(1)$	$\#V + \#E$
Dijkstra's	Priority Queue	$O(1)$	$O(\log(\#E))$	$O(\log(\#E))$	$(\#V + \#E) \log(\#E)$

```
SELECT expr1, SUM(expr2)
FROM dataframe
WHERE condition
GROUP BY expr1
```