

CSE 350

Advanced Data Structures

Topic 14: Column Stores

Example Queries

SELECT type, address FROM Permits
SELECT city, sum(price) FROM Permits
SELECT type, sum(price) FROM Permits WHERE city = 'Buffalo'

2 attrs/columns

3

I only need 2-3 cols
I need to load all data

↑ but I don't know which cols are needed
for any specific query

Tuple Linearization

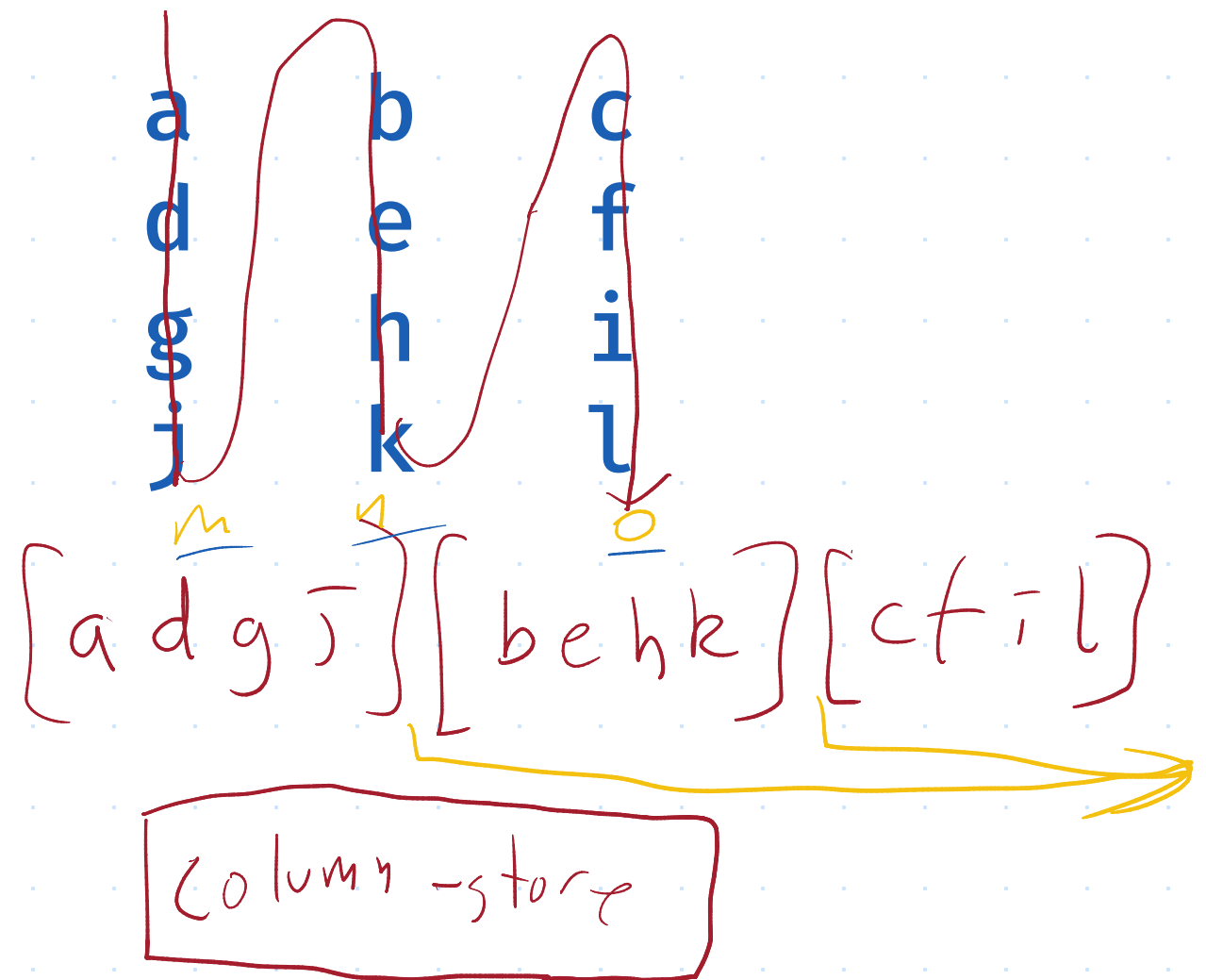
fields

records

a	b	c
d	e	f
g	h	i
j	k	l

[a b c][d e f][g h i] —

row-store



Column Storage (Naive)

Permits(id, type, address, city, price, ...)

Permits_id(rowid, id)

Permits_type(rowid, type)

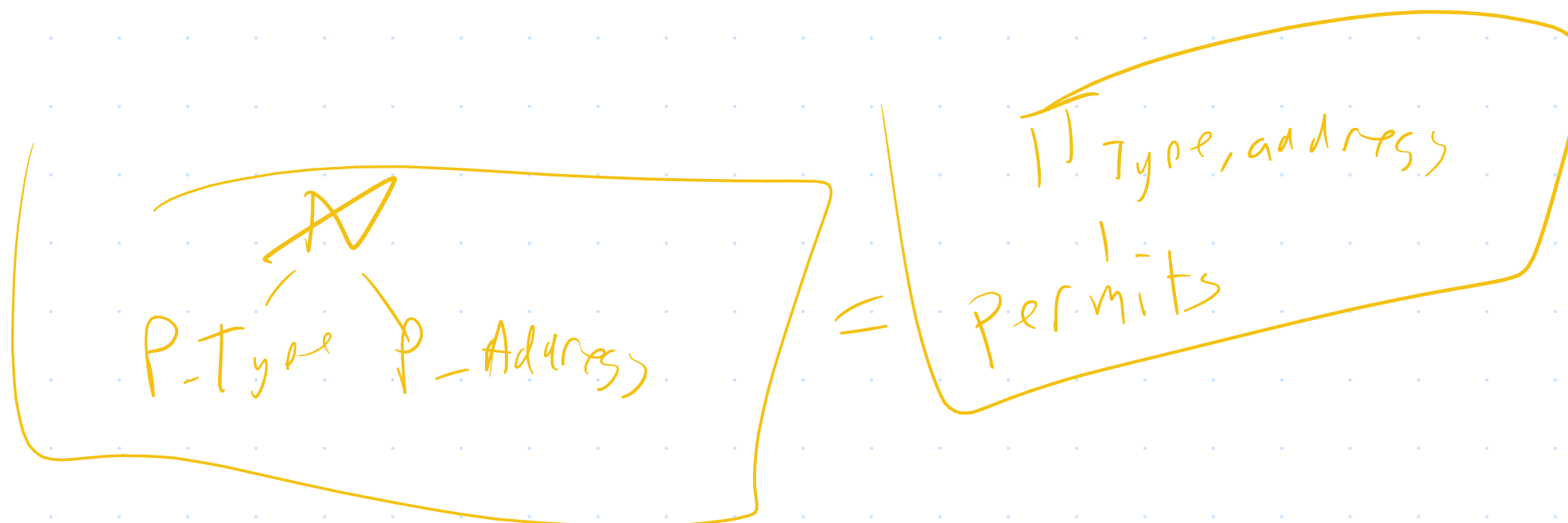
Permits_address(rowid, address)

— — —
— — —

Tuple Reconstruction Join

```
SELECT id, type, address, ...  
FROM permit-id  
JOIN permit-type ON rowid  
JOIN permit-address ON rowid  
-----
```

= Permits

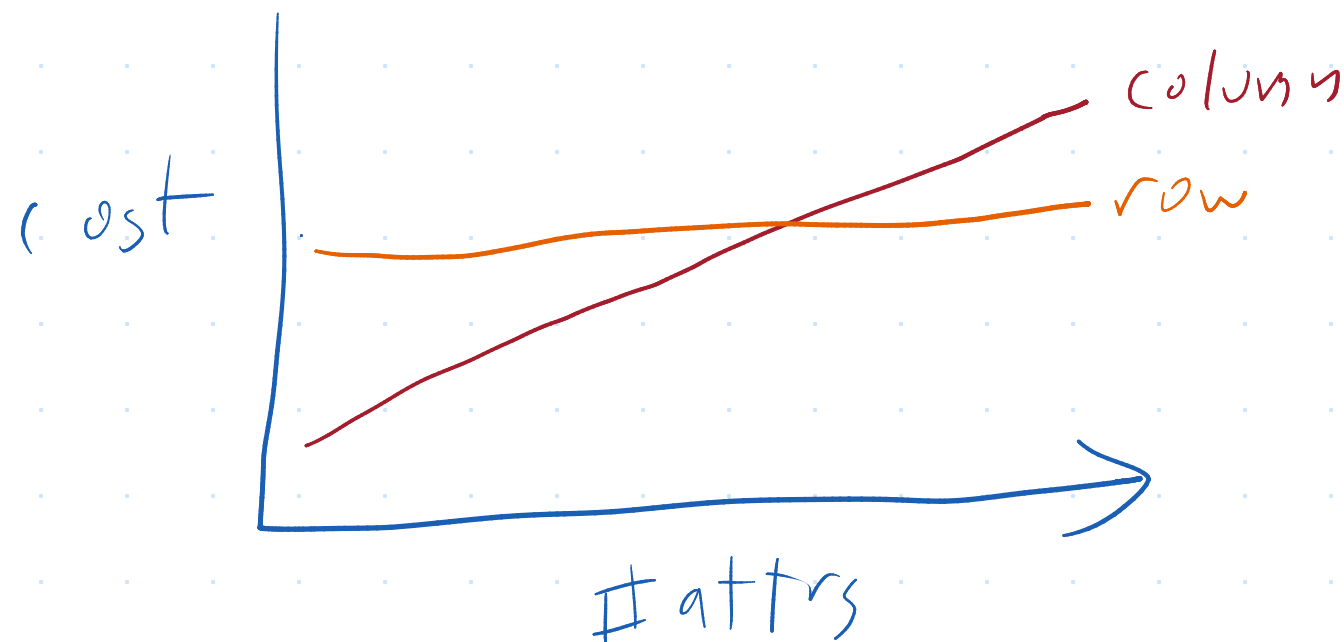


Con: Cost grows w/ # of attributes
↳ cost w/ all attrs is higher
↳ CPU cost goes up due to Join

More data stored

↳ need to store row id too

Update cost grows (one Io per attr)
↳ (worse w/ some optimizations)



Variation 1: Sort Order

Permits_Type sorted on type
Permits-Address sorted on address
Can have clustered indexes
Pro: on multiple attributes

Con

Permits-city sorted on city
Permits-price sorted on price

~~Permits-city~~
Permits-price

} tuple reconstruction going
gets much more expensive

↳ solution: Store 2 copies of each table

one sorted on attr
one sorted on rowid } more storage
(storage is cheap)
no more I/O needed

permits - city sorted on city
permits - price sorted on city } same sort order

↓
Correlated

↑
Join becomes a merge
(cheaper)

city							
price							

Extreme case: all attrs are correlated & rowid not needed

↳ city, address, zip all correlated
↳ rowid is implicitly position in the table

usual
case
mix
of both

↳ all other attrs stored in their own order
↳ rowid stored explicitly in the table

Implicit key/rowid

explicit rowid

Summary

can't
do
both
always

→ Sort order per column

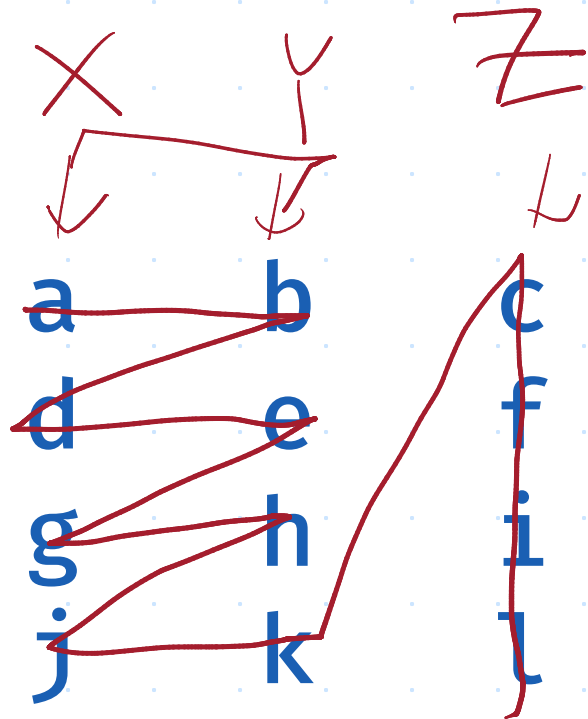
→ implicit keys (can't sort in multiple ways)

→ correlated columns

→ faster tuple reconstruction

Variation 2: Mix and Match

Column Groupings



[a b d e g h j k c f i l]

XY (rowid, X, Y)
Z (rowid, Z)

SELECT X, Y
FROM --

great
No tuple
reconstruction

SELECT X
FROM --

better than
row-oriented
but unnecessary
(read at Y)

Fractured Mirrors

a	b	c
d	e	f
g	h	i
j	k	l

m

n

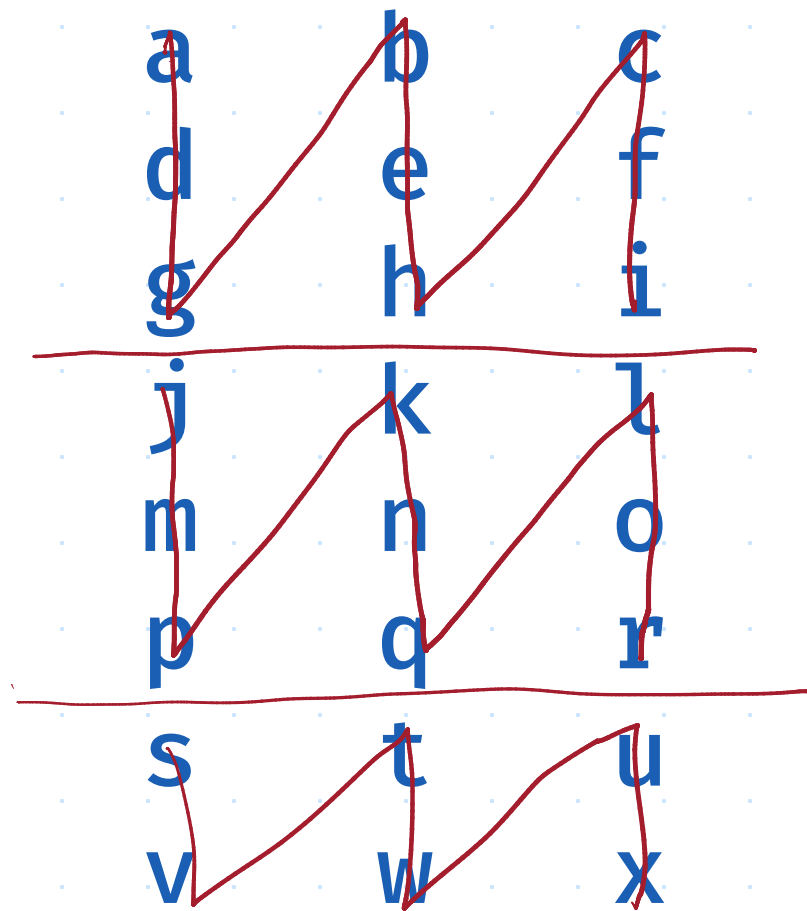
o

One copy (ground truth) in row oriented
one copy in column oriented

a	b	c
d	e	f
g	h	i
j	k	l

Used in many major DBs

PAX/Decomposed Storage

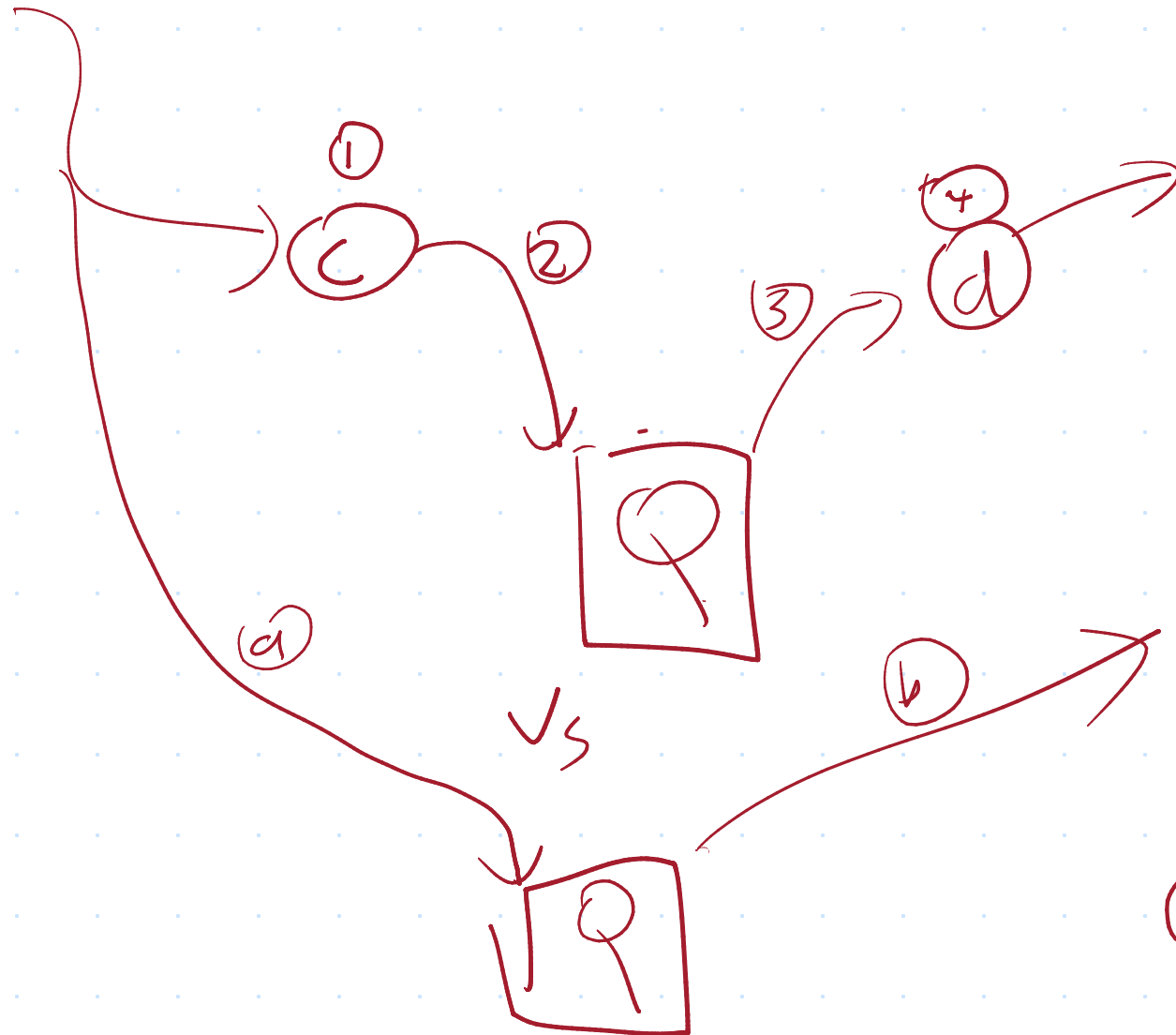


a d g b e h c f i j m p k n o l r s t u v w x

Compression

Permits_Type(rowid, type)

Permits_Address(rowid, address)



Goal

Idea: $(a) + (b) > (1+2) + (3) + (4)$

required $\Rightarrow (3) + (4)$

Con: decompression is not
free

pro = less IO (hopefully)

Compression

X

$Y = c(X)$

$X = d(Y)$

$|Y| \ll |X|$

Data

Compression

Decompression

Examples

Lempel-Ziv

Huffman Coding

DEFLATE (gzip)

Value Compression

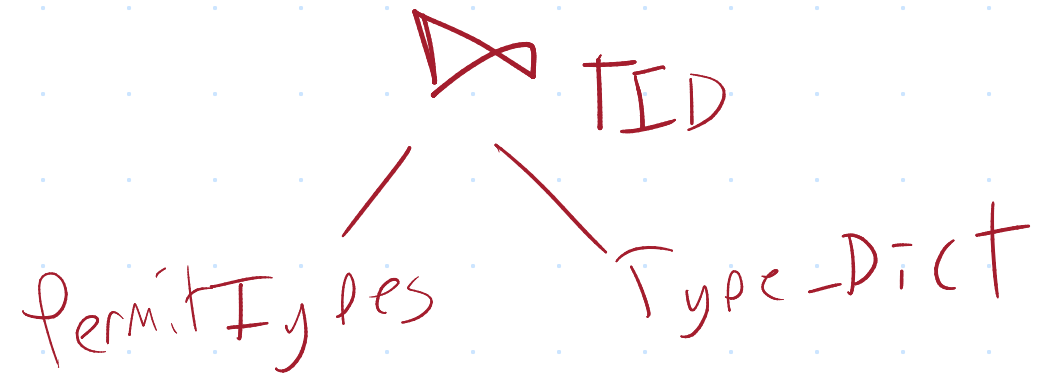
store $c(X)$

e.g. address is low entropy
↳ good compression ratio
 $|c(X)| \ll |X|$

Dictionary Encoding

<u>rowid</u>	<u>Type</u>	
0	USE	0
1	ELECTRICAL	1
2	ELECTRICAL	1
3	PLAN-COPY	2
4	ELECTRICAL	1
5	PLUMBING	3
6	ELECTRICAL	1
7	ELECTRICAL	1
8	ELECTRICAL	1
9	ASBESTOS	
10	ELECTRICAL	
11	HEATING	
12	ELECTRICAL	
13	PLAN-COPY	
14	HEATING	
15	PLUMBING	
16	ELECTRICAL	
17	REPAIR	
18	ELECTRICAL	
19	PLAN-COPY	

'0' ; 'U' 'S' 'E' 'n'
Type-Dict
Type TID
USE 0
ELECTRICAL 1
PLAN-COPY 2
;



Run-Length Encoding

<u>rowid</u>	<u>Type</u>	
0	USE	} 1
1	ELECTRICAL	
2	ELECTRICAL	} 2
3	PLAN-COPY	
4	ELECTRICAL	} 1
5	PLUMBING	
6	ELECTRICAL	} 3
7	ELECTRICAL	
8	ELECTRICAL	
9	ASBESTOS	
10	ELECTRICAL	
11	HEATING	
12	ELECTRICAL	
13	PLAN-COPY	
14	HEATING	
15	PLUMBING	
16	ELECTRICAL	
17	REPAIR	
18	ELECTRICAL	
19	PLAN-COPY	

(1, USE)
(2, Electrical)
(1, PLAN-COPY)
|
|
|

7-Bit Encoding

92 0101 1100
10399 0010 1000 1001 1111

01011100 1 byte

11010001 00011111